

Betting rules and information theory

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Outline

Simple betting in favorable games

The Central Limit Theorem

Optimal rules



The Game

Consider a simple game of chance, like tossing a coin or rolling some dice, in which you bet an amount x and if you win you get $2x$ and if you lose you lose the entire bet.

Let W_0 be the initial wealth so that if a fraction α is bet after the game you have

$$W_1 = \begin{cases} (1 - \alpha)W_0 + 2\alpha W_0 = (1 + \alpha)W_0 & \text{if you win} \\ (1 - \alpha)W_0 & \text{if you lose} \end{cases}$$

How much would you like to bet?



Classification of games

A game is *favorable* if the expected value of what you get is higher than what you bet.

A game is *fair* if the expected value of what you get is equal to what you bet.

A game is *unfair* if the expected value of what you get is lower than what you bet.



wealth maximization

Let the probability to win be p . The expected wealth after participation into the game can be computed

$$E[W_1] = (1 - \alpha)W_0 + p2\alpha W_0 = (1 - \alpha + 2\alpha p)W_0$$

If $p = 2/3$ then $E[W_1] = (1 + \alpha/3)W_0$ and to maximize the expected wealth one must chose $\alpha = 1$, so that $E[W_1] = 4/3W_0$.

The choice $\alpha = 1$ is made each time $2p - 1 > 0$, that is for any favorable game.



Now suppose that the game is repeated T times. If one plays with $\alpha = 1$ after T games it will have an expected wealth of

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With $T = 10$ the expected wealth is $E[W_T] \sim 17 W_0$ but the probability to have nothing is around 98%.



Consider the evolution of wealth in a sequence of repeated games

$$W_{t+1} = \begin{cases} (1 + \alpha)W_t & \text{if } \omega_t = 1 \\ (1 - \alpha)W_t & \text{if } \omega_t = 0 \end{cases}$$

where ω_t is a random variable taking value 1 for a win and 0 for a loss. The sequence of random variables $\omega = (\omega_1, \dots, \omega_T)$ is a *random process*.



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The final level of wealth depends on the number of 1's and 0's in the sequence ω .



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Example: $t = 2$ and $T = 3$. The number of sequence having 2 wins is 3 (the ways or “combinations” of selecting 2 elements in a set of 3) and the probability to obtain each sequence is $p^2 (1 - p)$.



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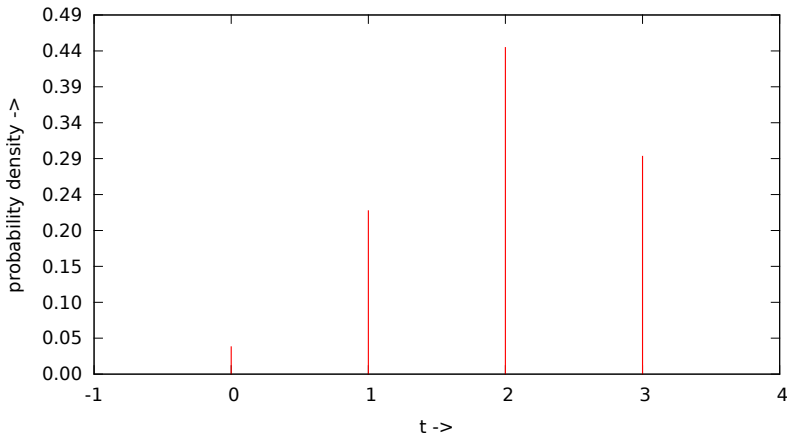
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In general how does the distribution of t looks like?



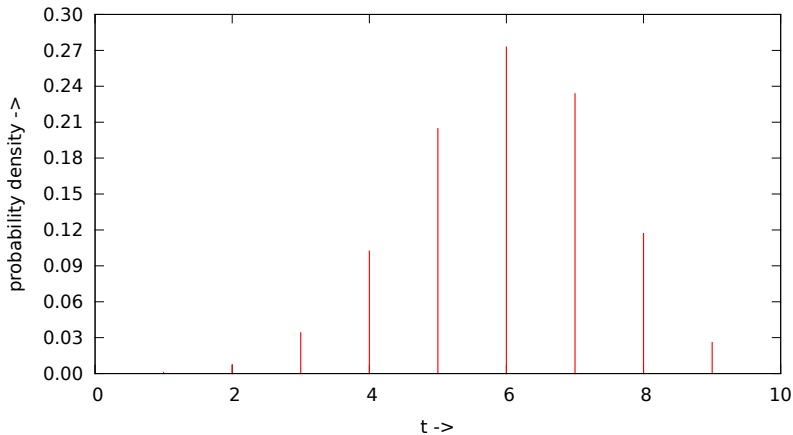
binomial with $T = 3$, $p = 2/3$



For $T = 3$ the possible values of t range from 0 to 3. Notice that $\text{Prob}[\text{wins} = 3]$ is larger than $\text{Prob}[\text{wins} = 0]$ because the probability of a win is larger than $1/2$.



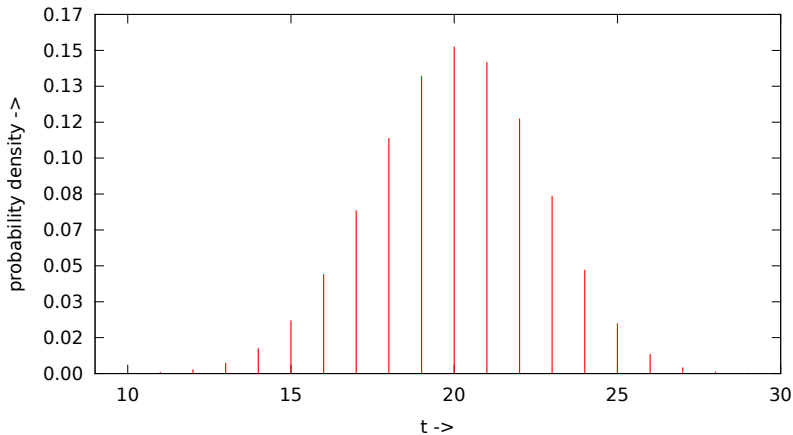
binomial with $T = 9$, $p = 2/3$



For $T = 9$ the possible values are from 0 to 9. The distribution is peaked around the mean and median value $p * T = 2/3 * 9 = 6$.



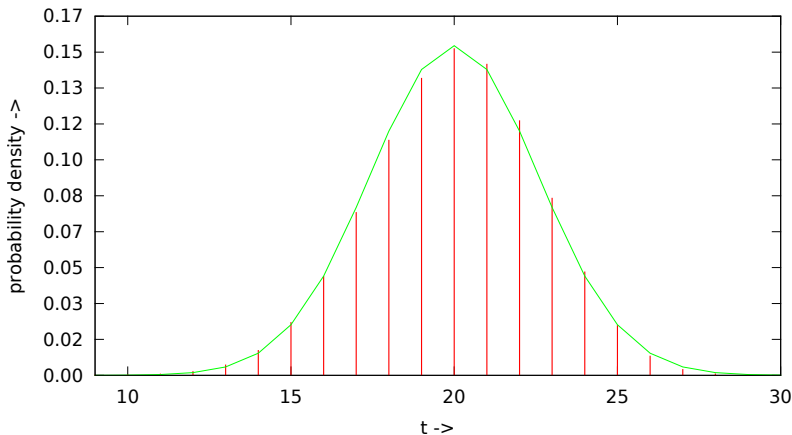
binomial with $T = 30$, $p = 2/3$



For $T = 30$ the distribution has already assumed a bell shape around $p * T = 2/3 * 30 = 20$.



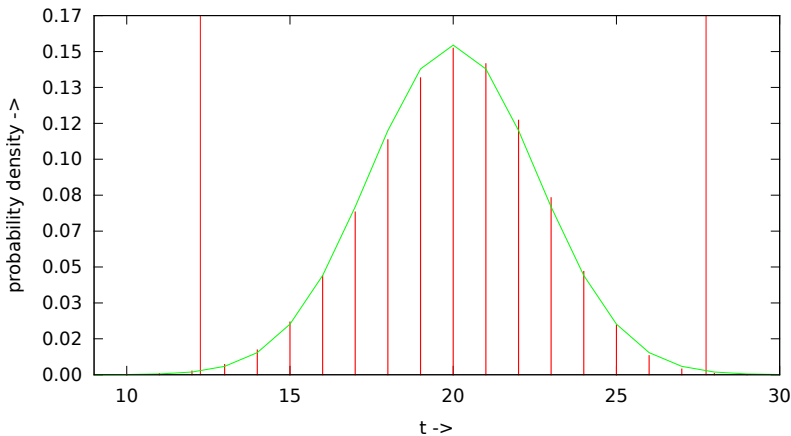
binomial and normal with $T = 30$, $p = 2/3$



The distribution is well approximated by a normal distribution with mean $\mu = pT$ and standard deviation $\sigma = \sqrt{p(1-p)T}$. In this case the mean is 20 and the standard deviation 2.58. This is the **Central Limit Theorem** at work.



binomial and normal with $T = 30$, $p = 2/3$

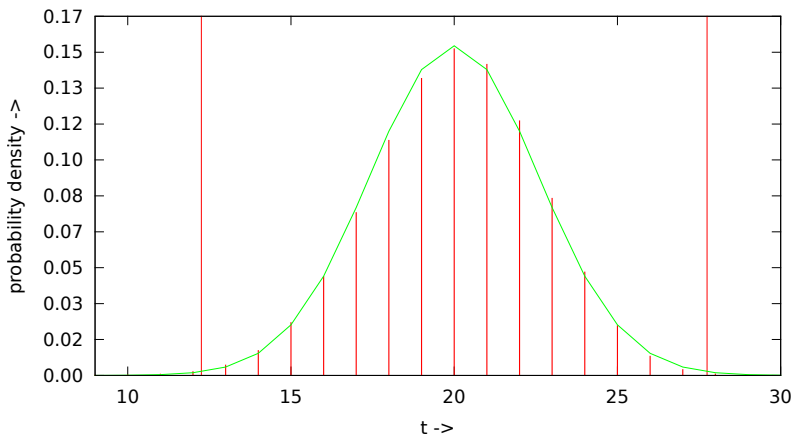


Moreover, due to the normal approximations, we can conclude that in the 99.73% of sequences the number of wins t will be in the interval

$$[\mu - 3\sigma, \mu + 3\sigma] .$$



binomial and normal with $T = 30$, $p = 2/3$



Different betting strategies work best in different cases: $\alpha = 0$ is the best if $t = 0$ and $\alpha = 1$ is the best if $t = T$. These are however unlikely events. It is more effective to have a strategy that works better around the modal value of the distribution of t .



The most probable level of wealth after T bets for a given α is $W_T^{\text{modal}} = (1 + \alpha)^{pT} (1 - \alpha)^{(1-p)T} W_0$. It grows exponentially with T .



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Substituting the minimal and maximal value of t for the interval derived before, in the 99.73% of sequences the wealth at time T satisfies

$$\frac{1}{T} \log \frac{W_T}{W_T^{\text{modal}}} \in \left[\frac{3\sqrt{p(1-p)}}{\sqrt{T}} \log \frac{1-\alpha}{1+\alpha}, \frac{3\sqrt{p(1-p)}}{\sqrt{T}} \log \frac{1+\alpha}{1-\alpha} \right].$$



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When T becomes large the interval is reduced to 0 so that, in probability

$$\lim_{T \rightarrow \infty} W_t = W_T^{\text{modal}}$$

which is just the **Law of Large Numbers**.



The best betting rule in the long run is the rule which better performs when the number of wins is equal to the modal value

$$\alpha^* = \arg \max \left\{ (1 + \alpha)^{pT} (1 - \alpha)^{(1-p)T} \right\}$$

removing the unnecessary T and taking the log

$$\alpha^* = \arg \max \{ p \log(1 + \alpha) + (1 - p) \log(1 - \alpha) \}$$



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Notice that the argument in the previous expression is just the expected log growth rate $E[\log(W_{t+1}/W_t)]$, thus α^* is the **log-optimal rule**.



In our case taking the derivative w.r.t. α one has

$$\frac{p}{1 - \alpha^*} - \frac{1 - p}{1 - \alpha^*} = 0$$

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According to the log-optimal rule: if the game is *favorable* ($p > 1/2$) bet a finite amount of money; if the game is *fair* ($p = 1/2$) bet nothing; if the game is *unfair* one wants to bet a negative amount, that is take the bet and commit herself to pay the possible payoff.

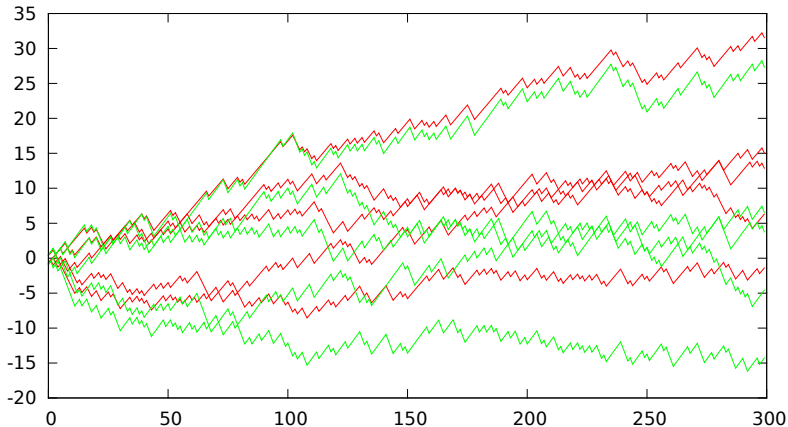


Wealth trajectories associated to different α 's will diverge. Set $w_0 = 1$ and imagine to play different sequences of games of length T with a given α and report the wealth levels $\log(W_t)$ on a graph. Then repeat the experiment with a different α .



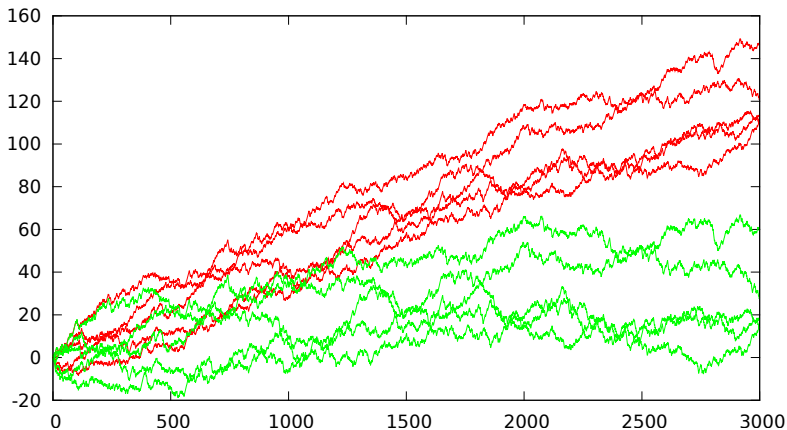
When the length of the sequence increases, the performances of the two strategies becomes more differentiated.

wealth trajectories for $\alpha = .5, .6$ and $p = 2/3$



On a longer time horizon the divergence is clear.

wealth trajectories for $\alpha = .5, .6$ and $p = 2/3$



expected wealth ratios

Starting with the same wealth at time t , consider the wealth at time $t + 1$ obtained with two different rules α and α'

$$W_{t+1}^{\alpha} = \begin{cases} (1 + \alpha)W_t & \text{if win} \\ (1 - \alpha)W_t & \text{if loss} \end{cases} \quad W_{t+1}^{\alpha'} = \begin{cases} (1 + \alpha')W_t & \text{if win} \\ (1 - \alpha')W_t & \text{if loss} \end{cases} .$$

then the expected ratio at time $t + 1$ is

$$E \left[\frac{W_{t+1}^{\alpha}}{W_{t+1}^{\alpha'}} \right] = p \frac{1 + \alpha}{1 + \alpha'} + (1 - p) \frac{1 - \alpha}{1 - \alpha'}$$



ratio-optimal rule

The question is whether there exists a betting rule α' such that for any α it is

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$$\mathbb{E} \left[\frac{W_{t+1}^{\alpha}}{W_{t+1}^{\alpha'}} \right] = \frac{p}{1 + \alpha'} + \frac{1 - p}{1 - \alpha'} + \alpha \left(\frac{p}{1 + \alpha'} - \frac{1 - p}{1 - \alpha'} \right)$$



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This can be made greater than 1 for some α if the expression inside parenthesis is different from zero. Thus the only solution is $\alpha' = 2p - 1 = \alpha^*$. **ratio-optimal = log-optimal**



absolute vs. relative

The log-optimality is related to the betting rules that *beats* all other rules: the probability to have a lower wealth than anyone else using a different rule tend to zero when the number of games increases.



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This idea is NOT related to the maximization of some function of the individual wealth, but is instead based on a notion of *relative* performance of one rule with respect to other rules.



Entropy

Given a random variable X taking N values with probabilities (p_1, \dots, p_N) its *entropy* is defined as

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The entropy (or negative information, or information loss)

- ▶ does not depend on the value taken by the random variable, but only on its probability distribution.
- ▶ is non-negative
- ▶ is maximal when the “information” of the variable is minimal, $p_i = 1/N$
- ▶ it is minimal when the random variable is maximally informative, i.e. there is a j such that $p_j = 1$



Relative entropy

Given two random variables X and Y taking values on a set of N events with probabilities (x_1, \dots, x_N) and (y_1, \dots, y_N) , the *relative entropy* or *Kullback-Leibler divergence* (sometimes distance) is

$$D(X|Y) = \sum_i x_i \log \frac{x_i}{y_i}$$

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The relative entropy

- ▶ is non-negative
- ▶ is asymmetric
- ▶ is minimal when X and Y coincides



entropy-optimal

The expression for the expected log growth rate can be rewritten as

$$E[\log(W_{t+1}/W_t)] = -p \log \frac{p}{(1+\alpha)/2} - (1-p) \log \frac{1-p}{(1-\alpha)/2} + p \log p + (1-p) \log(1-p) + \log(2)$$



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The betting rule can be seen as a random variable taking values on the two events *win* and *loss* with probabilities proportional to $1 + \alpha$ and $1 - \alpha$ respectively. Then

$$E[\log(W_{t+1}/W_t)] \sim -D(\{p, 1-p\} | \text{betting rule})$$

the relative entropy of the occurrences (win/loss) *given* the betting rule.



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Maximizing $E[\log(W_{t+1}/W_t)]$ is equivalent to minimize the relative entropy. **entropy-optimal = log-optimal.**



Summary

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Under some well specified conditions, the log-optimal rule is the rule that minimizes the conditional entropy

